

CSc 461/561
Multimedia Systems
Review on TCP/IP Networking

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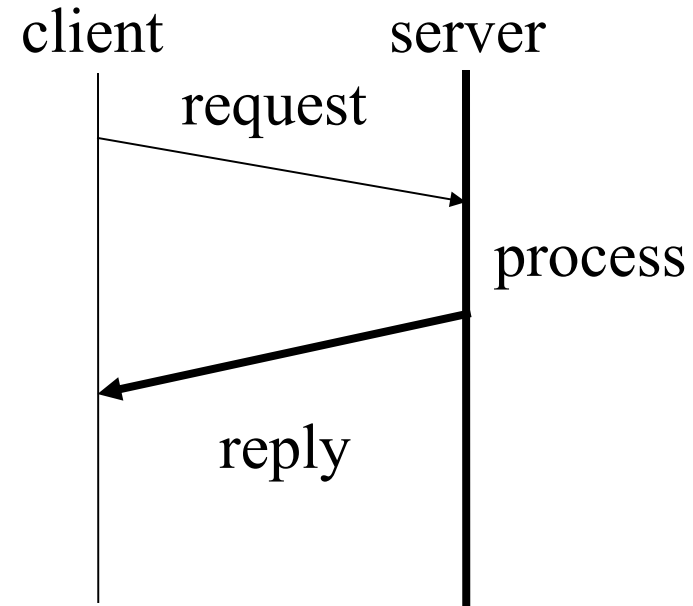
Application-oriented view

- (Old) Applications
 - remote login: e.g., telnet
 - file transfer: e.g., ftp
 - electronic mail: e.g., email
 - world-wide web: the Web!
- Requirements
 - move data from one location to another
 - elastic, error-free, in-sequence

* now: social networking, online audio/video streaming/gaming, etc

Client-server applications

- E.g., HTTP
 - HTTP client (browser)
 - **GET** /index.html HTTP/1.1
 - Host: www.example.com
 - (parameters)
 - HTTP server (Web server)
 - HTTP/1.1 **200 OK**
 - (metadata)
 - (data)



* now: ^{2/18/15}C/S-^{CSc 461/561}→CDN, P2P, cloud, hybrid, etc, but mostly ³based on HTTP

Protocols to support

- TCP/IP
 - the Internet Protocol Suite

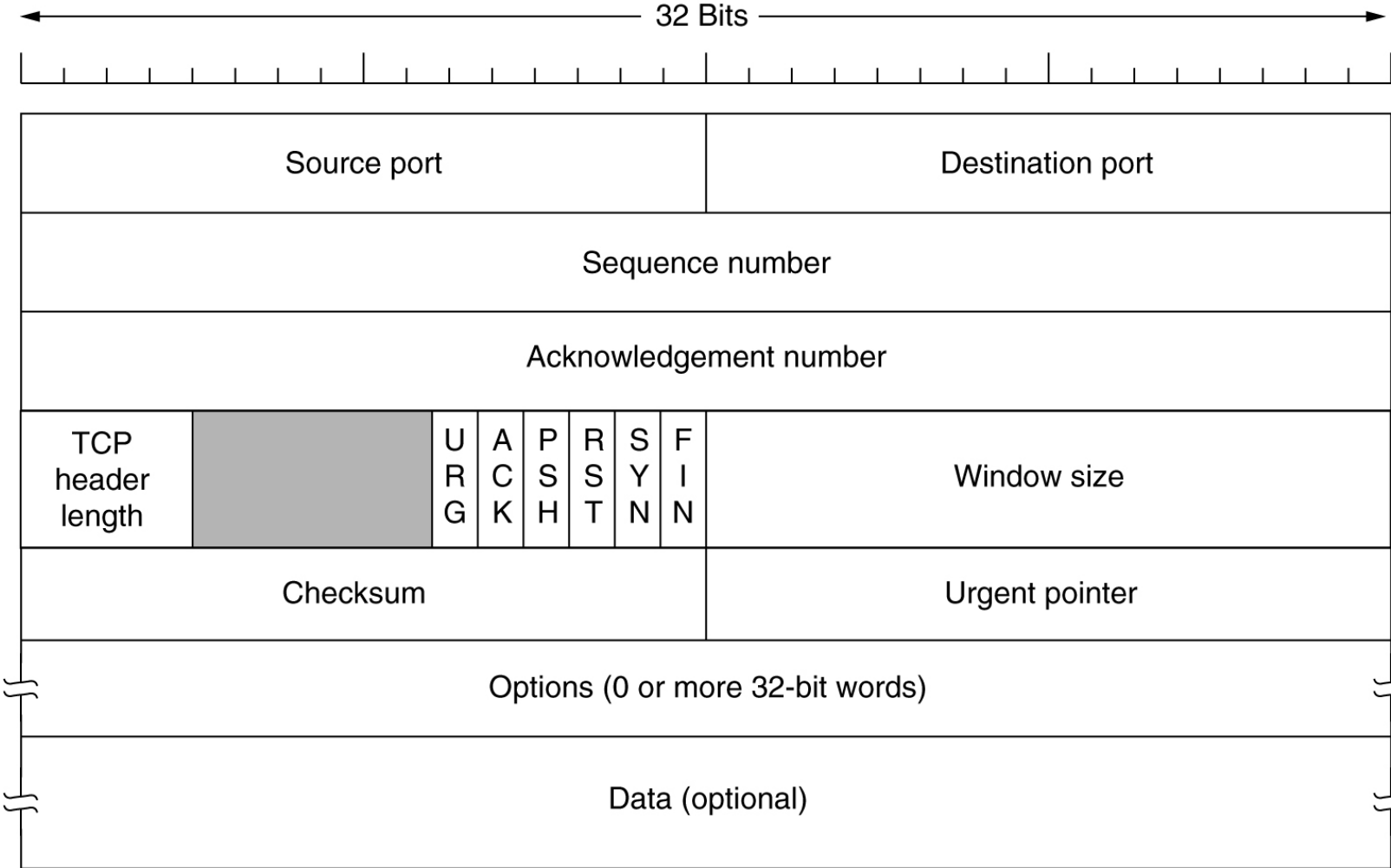
HTTP, etc	application
TCP/UDP	transport
IP	network
Ethernet, etc	link

- TCP offers
 - connection oriented
 - reliable, in-sequence, stream-like data transfer
- IP offers
 - addressing and routing; connectionless
 - IP packets may be lost, corrupted, duplicated, reordered

TCP

- Connection management
 - through packet handshake (SYN, FIN, ACK)
 - Multiplexing: port number
- Flow, error, congestion control
 - sequence number
 - acknowledgment number
 - window size
 - checksum

TCP header



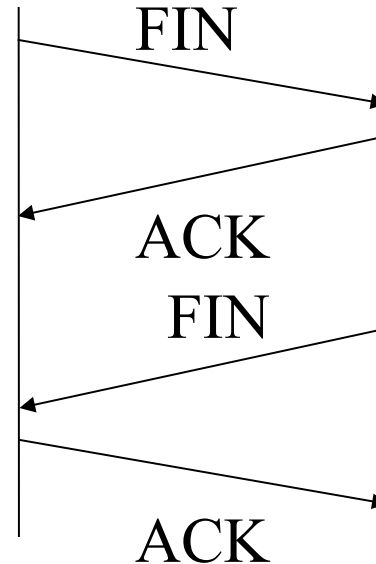
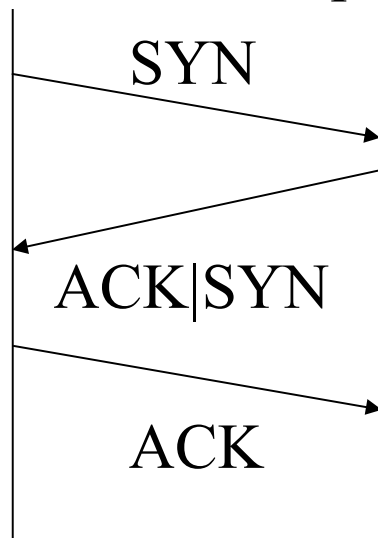
* 2/18/15 reserved/unused fields? Urgent Point?

Packet handshake

connection establishment

connection release

initiator responder



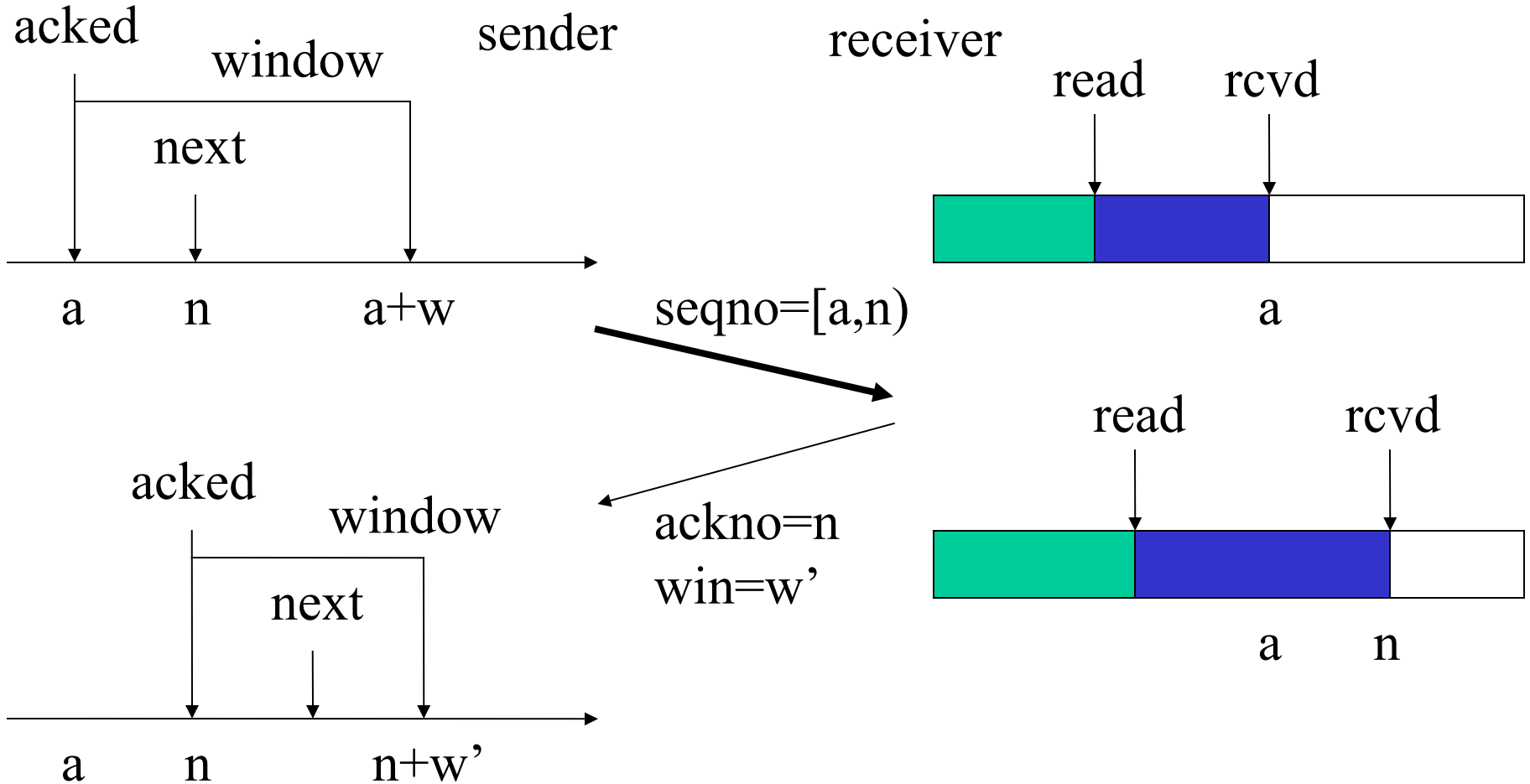
* 2/18/15 CSc.461/561
overhead for short connections and HTTP/0.9->1.0⁷->1.1, 2.0?

Flow control

- Purpose: pace sender and receiver
 - according to their buffer size
- Receiver's (advertised) window
 - available buffer space
- Sender's window
 - sliding window-based flow control
 - only send data within the window
 - sender's window $<$ receiver's window

* ^{2/18/15} other type of flow control: ^{CSc 461/561} rate-based, better for ⁸ MM applications?

Sliding window



Error control

- Error detection
 - receiver: sequence number, TCP checksum
 - sender: timeout
- Error notification
 - receiver=>sender: duplicate acknowledgment
- Error recovery
 - sender: end-to-end retransmission

Go-back-N retransmission

- Sender: send packet 1, 2, 3, 4, 5, 6
- Receiver: receive packet 1, 2, 3, 5, 6
 - cumulatively acknowledge up to packet 3
- Sender: send packet 7, 8, 9
 - timeout for packet 4; retransmit packet 4
- Receiver
 - cumulatively acknowledge up to packet 9
- Sender: send packet 10, 11, 12, ...

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* data may come too later or in duplication?

Congestion control

- Was not there when TCP was first designed
- Added to TCP since late 80s
 - heavily coupled with flow/error control
 - heavily explored research topics in decades!
- Purpose: pace sender and network
 - competing flows
 - overload network (e.g., output queue)
 - cause packet losses when queues overflow

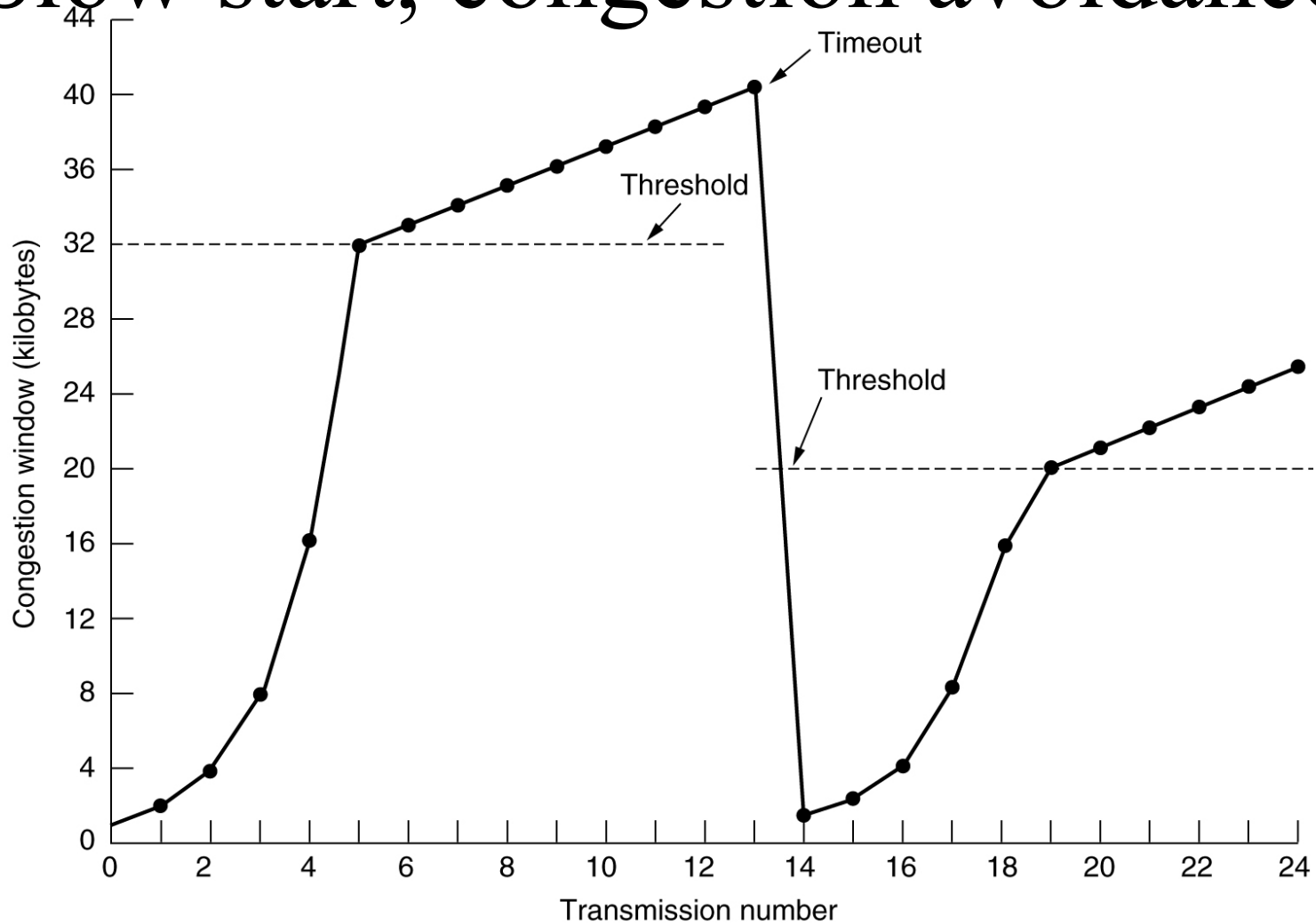
Congestion window

- Slow start
 - start with an initial congestion window (cwnd)
 - usually initial cwnd = 1 full-size packet
 - double window size every round-trip time until cwnd is above slow start threshold (ssthresh)
- Congestion avoidance
 - increase window size linearly

Congestion window: more

- Back-off
 - when timeout occurs, assume packet loss happened [error control]
 - $ssthresh = 0.5 * current_cwnd$
 - restart with the initial cwnd
 - retransmit with doubled timer [error control]

Slow start, congestion avoidance



*2/18/15
consequence on multimedia traffic?

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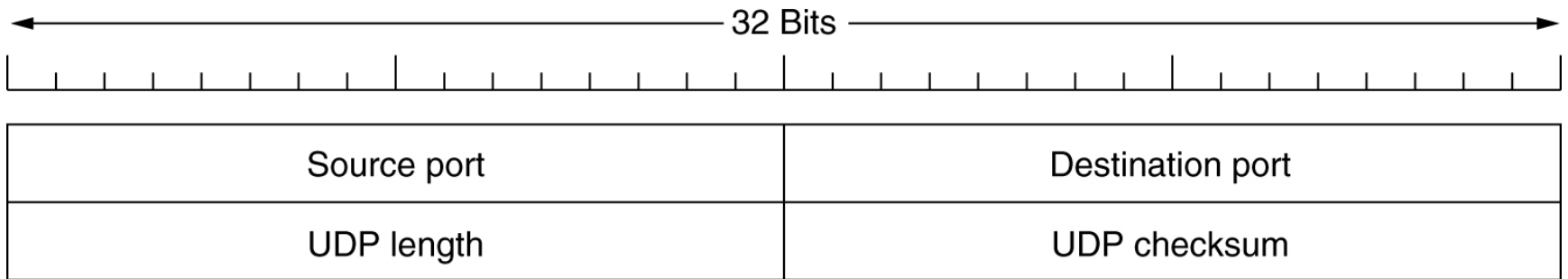
Congestion control: more

- Fast retransmit
 - retransmit with 3 duplicate acknowledgments
 - slow start threshold (ssthresh) to be a half of current congestion window (cwnd)
 - restart with the initial congestion window
- Fast recovery
 - similar to Fast retransmit
 - but restart with $cwnd = ssthresh$

UDP

- Why TCP is not enough?
 - sometimes TCP is an overkill
 - e.g., loss-tolerant, delay-sensitive applications
- UDP offers
 - connectionless, datagram-like data transfer
 - no reliability, in-sequence guarantee
 - *flexibility* to plug-in flow/error/congestion etc in application layers

UDP header



This lecture

- A quick review on
 - Internet Protocol Suite
 - TCP
 - UDP
- Explore further
 - Why TCP header has no “TCP length” field such as “UDP length” in UDP header?

Next lecture

- IP
- Link: wired and wireless
- Why multimedia networking is different, and difficulty?