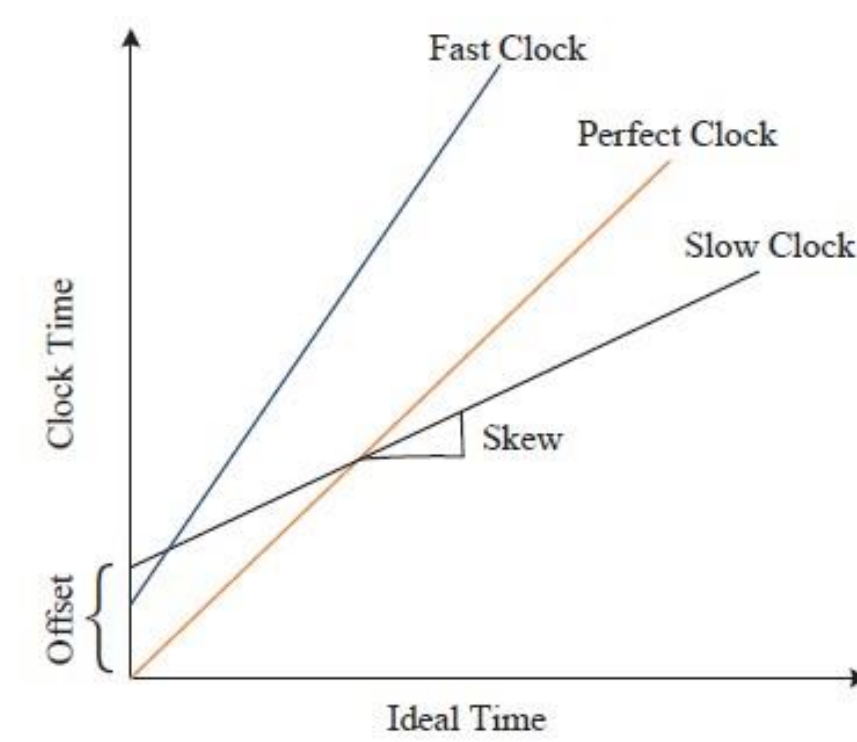


Background

What is time synchronization?

- Common clock skew (clock speed)
- Common clock offset (instantaneous clock difference)
- Getting all devices in network to the same time (clock offset) at exactly the same rate (clock skew)



Why do we need time synchronization in WSNs?

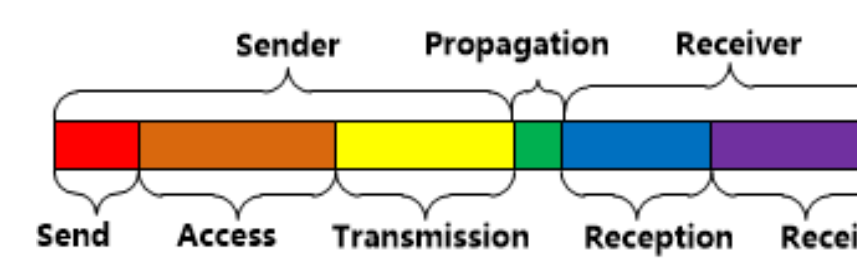
- Fundamental requirement of applications
- Precondition of sensors coordination
- Save sensors energy

How is time-synch in WSNs different from traditional networks?

- **Energy Utilization**
- Single hop vs. **multi hop**
- Infrastructure-Supported vs. **Ad-hoc**
- Static topology vs. **Dynamic Topology**
- Connected vs. **Disconnected**
- **Dynamic time sync. requirements**, depending on the application

What are the challenges in WSNs?

- Difference noises are inevitable in network
- Time-varying clock speed
- Distributed requirement
-
- i.e., simple problem but hard to be solved



Main Approaches

Maximum consensus-based approach

- maximum consensus-based for fast convergence
- i.e., always select the fastest clock as the reference clock

Stochastic approximation approach

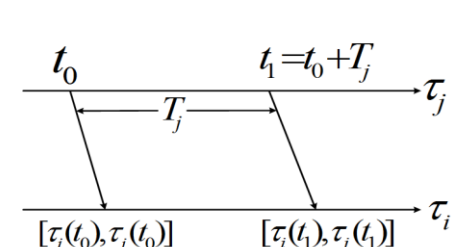
- make an average of all time of estimation to achieve accurate relative skew

Hops-based weight design and reference node number mark

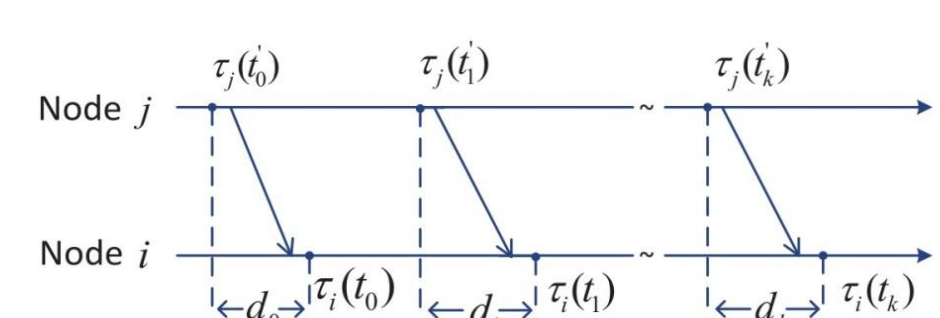
- design hops-based weight and mark the reference node for preventing clock drift

Relative Clock

Noise free

$$a_{ij} = \frac{\tau_j(t_i) - \tau_j(t_0)}{\tau_i(t_i) - \tau_i(t_0)} = \frac{a_j}{a_i}$$


Random noise

$$a_{ij}(k) = \frac{\tau_j(t'_k) - \tau_j(t'_{k-1}) + (k-1)a_j(k-1)}{\tau_i(t_k) - \tau_i(t_{k-1})}$$


Algorithms

MTS

$$\begin{cases} \hat{a}_i(k+1) = a_{ij}(k)\hat{a}_j(k) \\ \hat{b}_i(k+1) = \hat{a}_j(k)\tau_j(k) + \hat{b}_j(k) - \hat{a}_i(k+1)\tau_i(k) \end{cases}$$

$$\hat{b}_i(k+1) = \max_{x=i,j} [\hat{a}_x(k)\tau_x(k) + \hat{b}_x(k)] - \hat{a}_i(k+1)\tau_i(k)$$

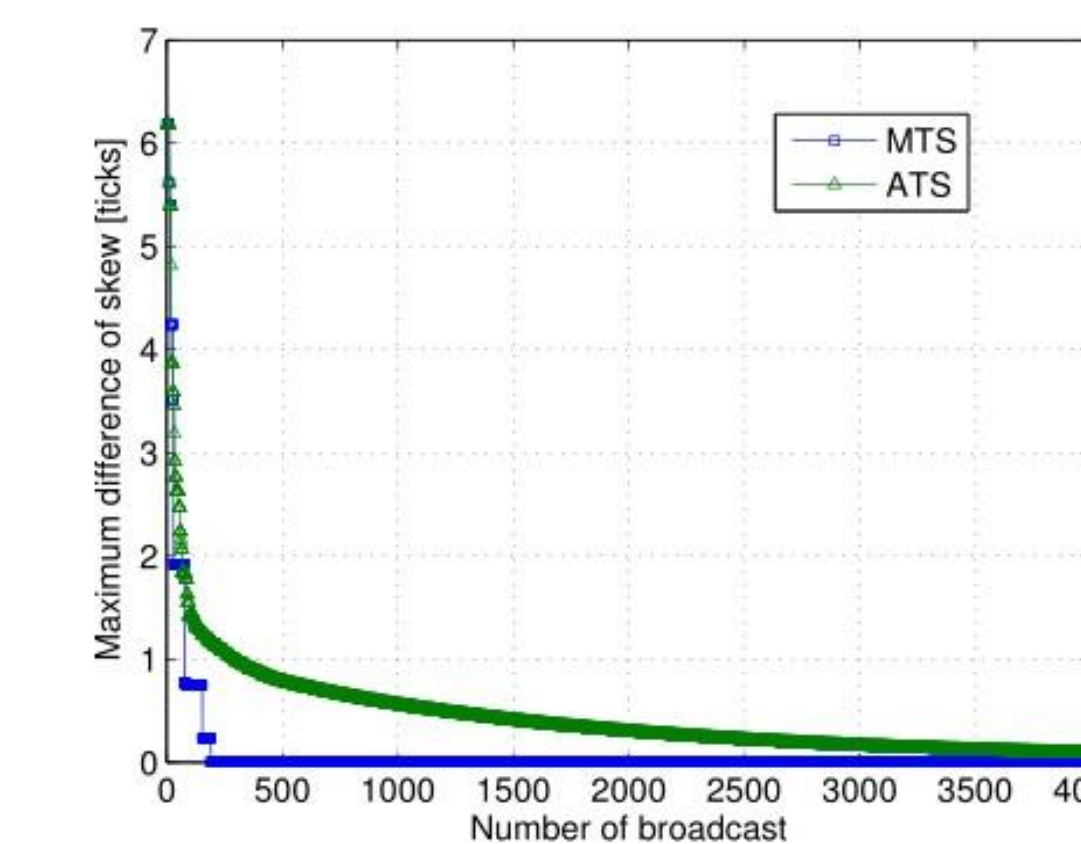
WMTS

$$w_i = w_j + 1 \quad r_i = r_j$$

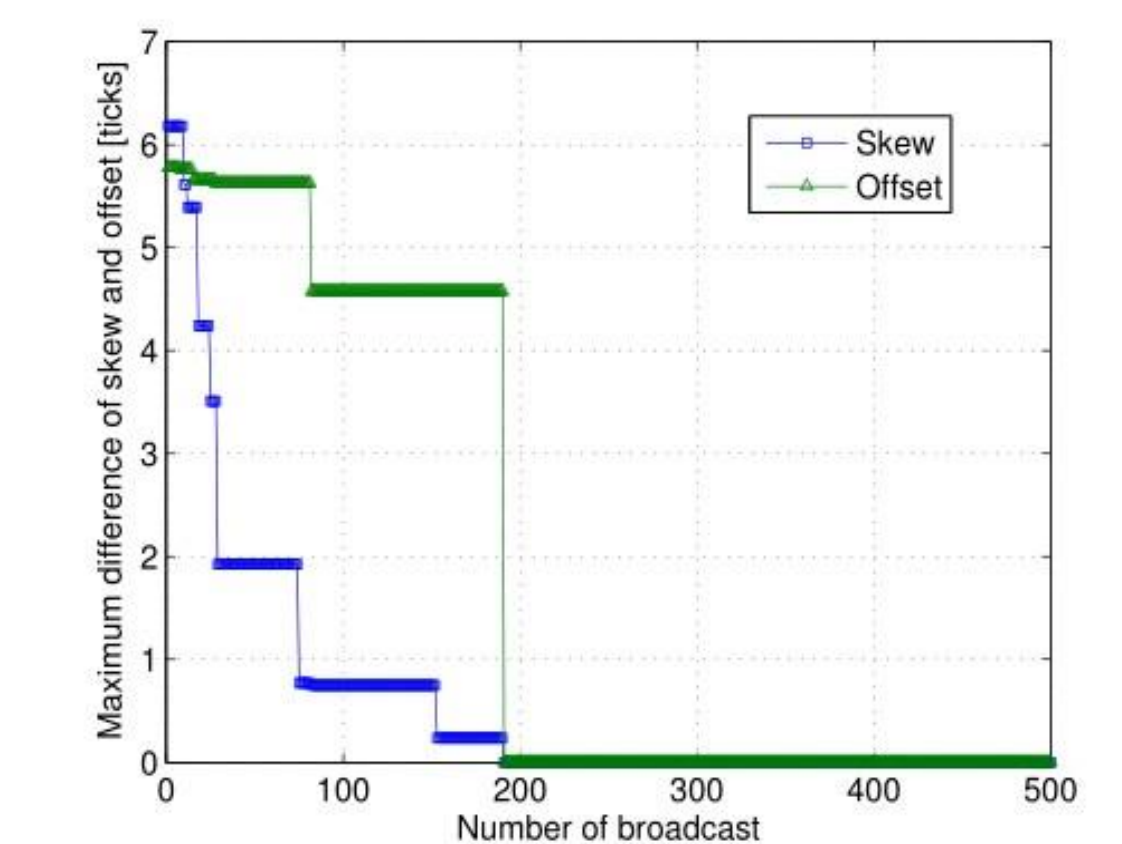
$$\begin{cases} \hat{a}_i(k+1) = a_{ij}(k)\hat{a}_j(k) \\ \hat{b}_i(k+1) = \hat{a}_j(k)\tau_j(k) + \hat{b}_j(k) - \hat{a}_i(k+1)\tau_i(k) \\ \hat{b}_i(k+1) = L_j(k) - \hat{a}_i(k)\tau_i(k) \end{cases}$$

Evaluation

Noise free

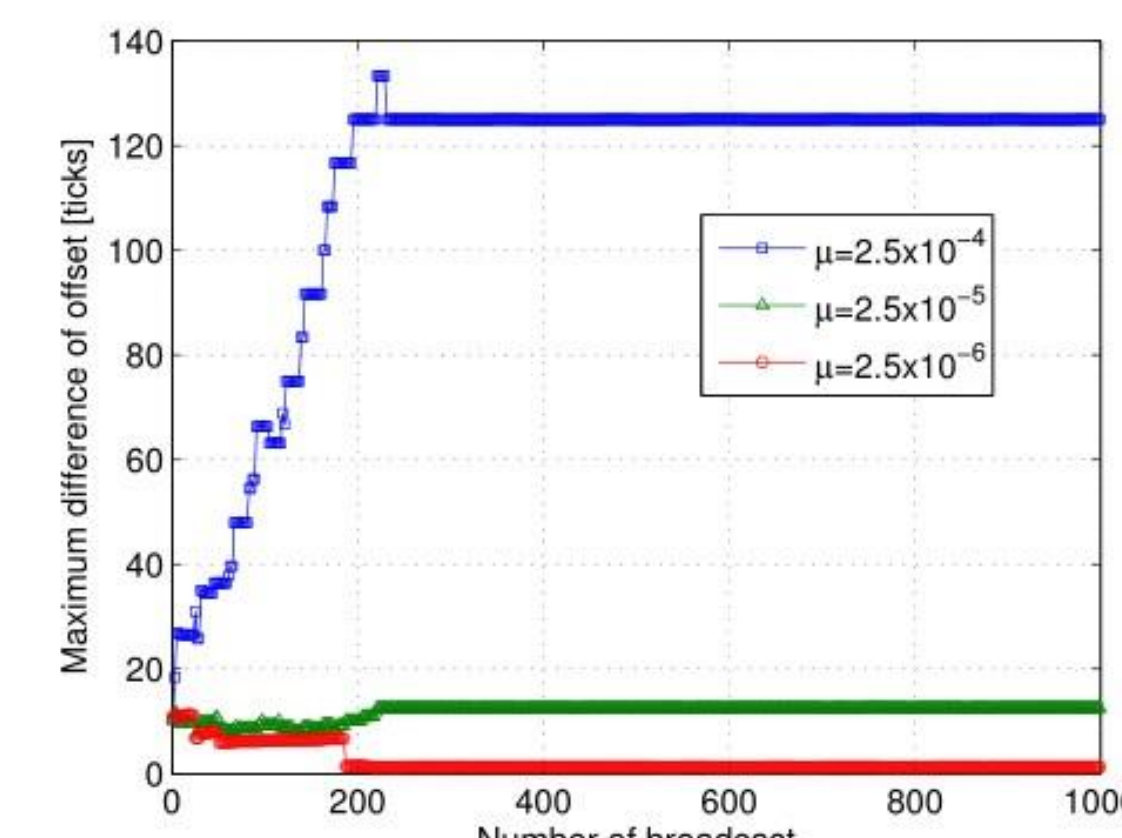
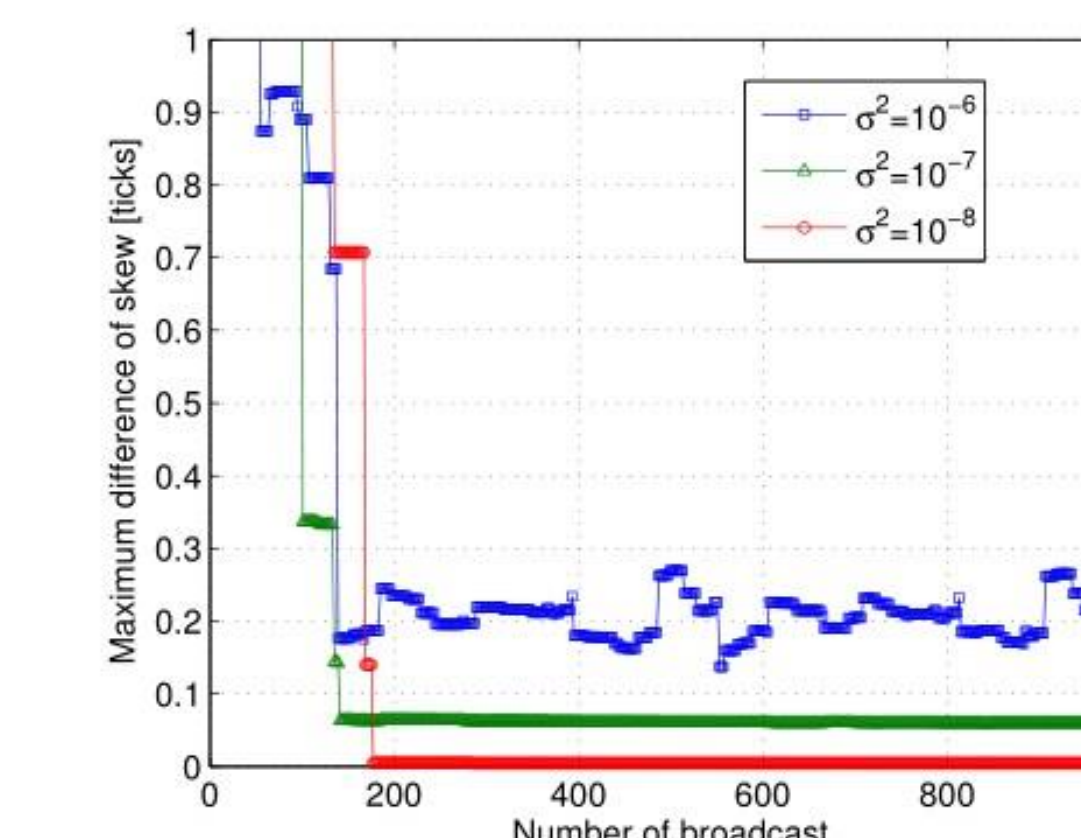
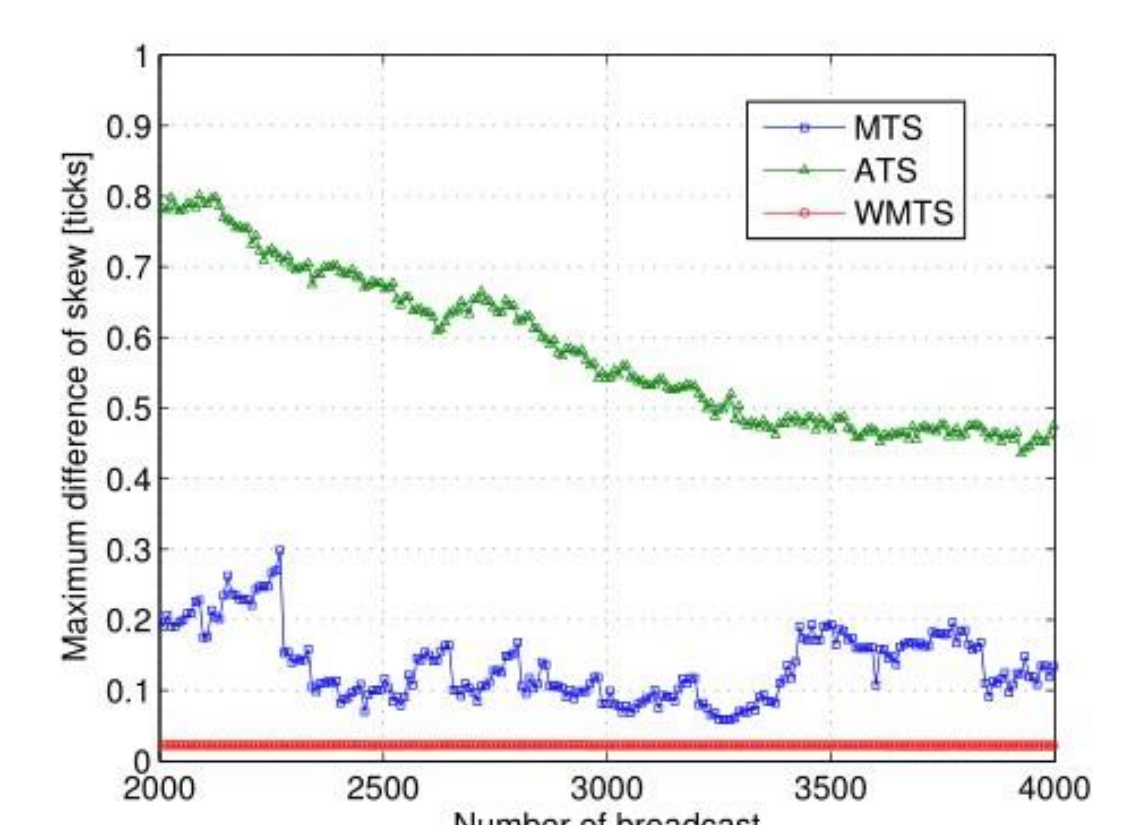
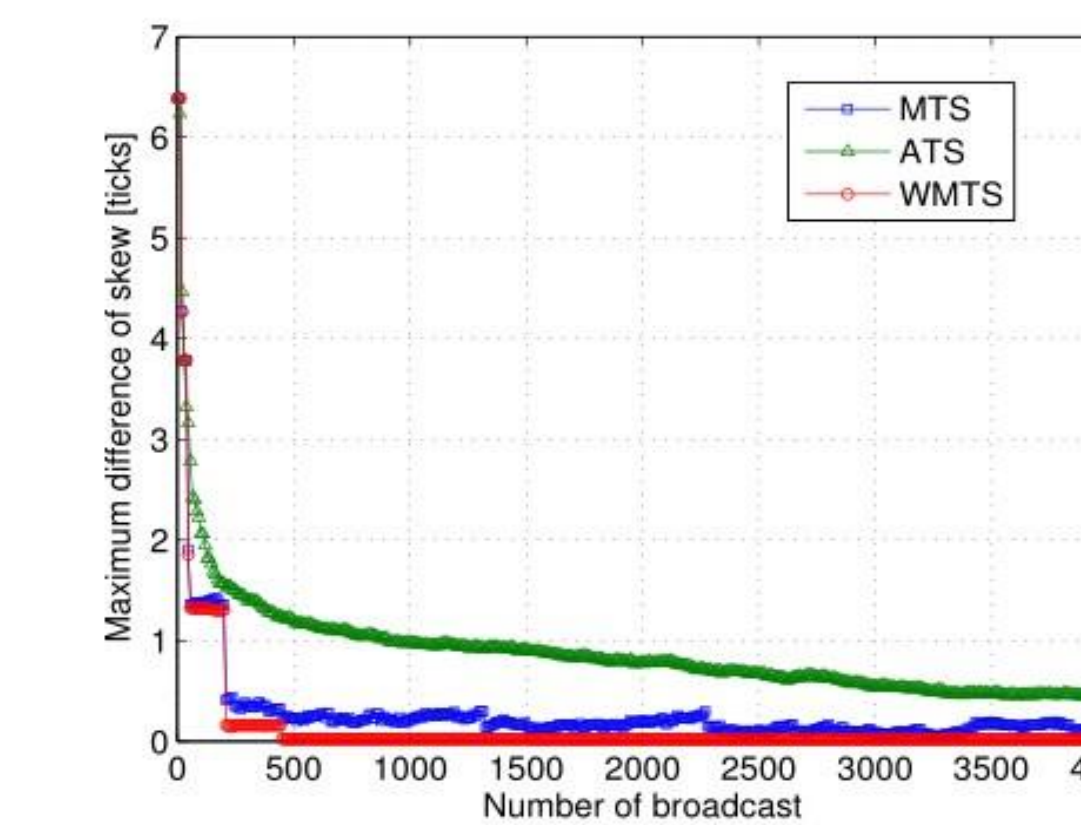


much faster convergence speed



compensate simultaneously

Random noise



higher synchronization accuracy and robust against to random noise

Motivation

Disadvantages for traditional protocol [1-2]

- Need a root or a reference node
- Tree-topology based
- Inaccurate synchronization
- Lack theoretical support

Advantages of consensus-based protocol [3-4]

- Fully distributed
- Strong robustness and scalability
- Highly accurate synchronization
- Compensate both skew and offset
- Exponential convergence speed

Improvements of our protocol [5-7]

- Finite-time convergence
- Solve the communication delay
- Lower complexity
- Compensate simultaneously
- Theoretical support

Clock Model

Hardware clock: linear model

$$\tau_i(t) = a_i t + b_i$$

Relative clock between nodes

$$\tau_i(t) = \frac{a_i}{a_j} \tau_j(t) + b_i - \frac{a_i}{a_j} b_j$$

Software clock: linear function

$$L_i(t) = \hat{a}_i \tau_i(t) + \hat{b}_i$$

Objective

Designed synchronization protocol, such that

$$\begin{cases} \lim_{k \rightarrow \infty} \hat{a}_i(k) a_i = a_c \\ \lim_{k \rightarrow \infty} [\hat{a}_i(k) b_i + \hat{b}_i(k)] = b_c \end{cases}$$

i.e., have the same logical clock

Main Results

Convergence of MTS

Theorem 1: asymptotic convergence of MTS

Considering a connected network, by MTS, the skew and offset converge and satisfy

$$\begin{cases} \lim_{k \rightarrow \infty} \hat{a}_i(k) a_i = a_{\max} \\ \lim_{k \rightarrow \infty} \hat{a}_i(k) b_i + \hat{b}_i(k) = b_{\max} \end{cases}$$

Theorem 2: finite-time convergence of MTS

Considering a connected network, by using MTS, the convergence time satisfy

$$T_{con} \leq B(N-1)$$

Convergence of WMTS

Theorem 3: mean-square convergence of relative skew estimation

$$E\{a_{ij}(k)\} = \frac{a_i}{a_j}, k \in N^+$$

Theorem 4: finite-time converge in expectation of WMTS

Considering a connected network, by WMTS, the converge in expectation

$$E\{\hat{a}_i(t) a_i\} = a_{\max}, i \in V, t \geq B(n-1) \quad E\{L_i(t)\} = a_{\max} t + b_{\max} - w_i(t) a_{\max} \mu, i \in V, t \geq B(n-1)$$

Theorem 5: mean-square convergence of WMTS

Considering a fixed connected network, by WMTS, we have

$$E\{\hat{a}_i(k) a_i - \hat{a}_j(k) a_j\} = 0 \quad \lim_{k \rightarrow \infty} \text{Var}\{\hat{a}_i(k) a_i - \hat{a}_j(k) a_j\} = 0$$

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